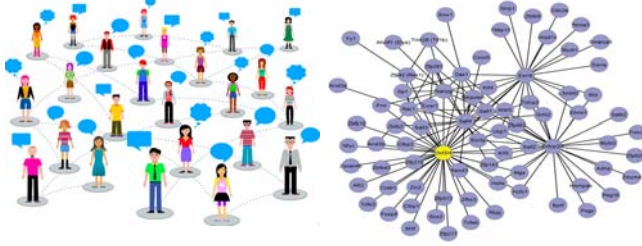


Streaming Link Prediction on Dynamic Graph Streams



Introduction

Real-world **information networks** are **dynamic** and **of massive volume**. Noteworthy examples include social networks, PPI networks, communication networks, and the Web.



Link Prediction

Given a snapshot of an information network at time t , we seek to accurately predict the edges that will be added to the network during the interval from time t to a given future time t_0

Graph Stream Model

An information network can be formalized as a graph stream that receives a sequence of edges of the form $\langle EId; i, j \rangle$. At any given moment in time t , the edges in the stream imply a graph $G(t) = (N(t), E(t))$, where $N(t)$ is the set of nodes and $A(t)$ is the set of *distinct* edges at time t

- **Dynamics**: t can be infinitively large
- **Massiveness**: $G(t)$ is too large to be safely stored even on disk for effective analysis



Can we predict potential links on massive, dynamic graph streams accurately and efficiently?

Main Idea

We consider proposing and re-examining high-quality **approximations** to the state-of-the-art link prediction metrics on the massive graph stream scenario:

1. Common neighbor:

$$|\tau(i, t) \cap \tau(j, t)|$$

2. Jaccard coefficient

$$\frac{|\tau(i, t) \cap \tau(j, t)|}{|\tau(i, t) \cup \tau(j, t)|}$$

3. Adamic-Ada

$$\sum_{k \in \tau(i, t) \cap \tau(j, t)} (1/\log(|\tau(k, t)|))$$

Algorithms

1. Minhashing based approximation

Consider a streaming algorithm maintaining values of each node i , its *minimum adjacent hash value*, $v(i)$, and *minimum adjacent hash index*, $I(i)$

Theorem: The probability that $I(i) = I(j)$ is exactly equal to the Jaccard coefficient between nodes i and j

2. Node-biased sampling

A **reservoir** of budget L is associated with each node to dynamically sample L incident edges of each node

Theorem: Let (i) and (j) be the fraction of nodes incident on i and j respectively, which are sampled. Then, the total number of common neighbors C_{ij} , between nodes i and j , can be estimated

$$C_{ij} = \frac{|S(i) \cap S(j)|}{\min\{\eta(i), \eta(j)\}}$$

Experiments

Real-world Datasets

- **DBLP**: a streaming co-authorship graph comprising 1,954,776 author-pairs
- **Amazon Co-purchasing Network**: a product co-purchasing network comprising 410,236 nodes and 3,356,824 edges

Evaluation Metrics

- Link prediction accuracy
- Link prediction cost

Experimental Results

